

IN THE CLAIMS

Please amend the claims as follows:

Claims 1 -7 (Cancelled).

Claim 8 (Previously Presented): An authentication process involving a first device, which possesses a public key v and a secret key s , the public and secret keys being related by an operation modulo n , where n is an integer, the modulus n being specific to the first device, and a second device, which knows the public key v , the first and second entities being provided with means to exchange zero-knowledge information and to carry out cryptographic calculations on the zero-knowledge information, calculations being carried out modulo n wherein in the process the modulo n operation is of $v=s^{-1} \pmod{n}$, t being a parameter and in that the modulo n calculations are performed according to the "Chinese remainders" method and in that the modulus n is the product of two primes of similar size.

Claim 9 (Previously Presented): A process according to claim 8, wherein the information exchanges are of zero-knowledge and wherein the cryptographic calculations are completed as follows:

the first device selects at least one integer r at random ranging between 1 and $n-1$ and calculates at least one parameter x equal to $r^t \pmod{n}$, then at least one number c that is at least one function of the at least one of a parameter and a message and sends the at least one number c to the second device;

the second device receives the at least number c , selects at least one number e at random, and sends the at least one number e to the first device;

the first device receives the at least one number e , carries out at least one calculation using the at least one number e and the secret key s , the result of the at least one calculation yielding at least one answer y , and sends the at least one answer y to the second device.

the second device receives the at least one answer y , carries out one calculation using the public key v and the modulus n , and checks with a modulo n operation that the result of the one calculation is coherent with the received at least one number c .

Claim 10 (Previously Presented): A process according to Claim 9, wherein a size of the number n , expressed in number of bits, is less than 1,000.

Claim 11 (Previously Presented): A process according to Claim 10, wherein a size of the number n is between 700 and 800.

Claim 12 (Cancelled).

Claim 13 (Currently Amended): A message signature process configured for a device provided with a public key v and a secret key s , the public and private keys being related by a modulo n calculation, where n is an integer, which is specific to the device, the process utilizing means configured to calculate at least one number c that is a function of a message M to be signed, configured to calculate at least one number y that is a function of the secret key s , and configured to transmit the numbers y and c that are the signature of the message and the message M , wherein the modulo n operation is $v=s-t \pmod n$, t being a parameter wherein the modulo n calculations are performed according to the "Chinese remainders" method and in that the modulus n is the product of two primes of similar size.

Claim 14 (Previously Presented): A message signature process according to claim 13, wherein the device selects an integer r at random between 1 and $n-1$, calculates a parameter x equal to $rt \pmod n$, calculates at least one number e that is a function of parameter x and the message M to be signed, calculates the at least one number y using its secret key s , said at least one number y being a function of numbers r and e , and transmits the numbers c and y as the signature.